

Class Notes - CS482 - Chapter 2

1 A Student Registration System

A project, that starts with a statement of the following objectives:

1. Authenticate themselves as users of the systems
2. Register and deregister for courses
3. Obtain report on a particular student's status
4. Maintain information about students and courses
5. Enter final grades for courses that a student has completed.

We will have the first specification similar to this for our project.

2 Relational Databases

Mainly interested in **relational model**.

Relation and tuples. Databases are stored in *tables*. Each row in a table represents an object or an entity of interest. Each column in the table describes the object in a particular way. Columns are *attributes* and each has its *domain*.

Example 1 *The student registration system should contain a table about students with the attributes like Id, Name, Address, Status. Each row represents a student. Each column describes the student in a way. The current information might look as follows:*

<i>Id</i>	<i>Name</i>	<i>Address</i>	<i>Status</i>
111111111	John Doe	123 Main St.	Freshman
666666666	Joseph Public	666 Hollow Rd.	Sophomore
111223344	Mary Smith	1 Lake St.	Freshman
987654321	Bart Simpson	Fox 5 TV	Senior
023456789	Homer Simpson	Fox 5 TV	Senior
123454321	Joe Blow	6 Yart Ct.	Junior

The STUDENT Table

Relation. A relation is a set of *tuples*. A relation could be viewed as a predicate with variables. The set of tuples defines the predicate (if a tuple belongs to the relation then it is true; otherwise, it is false)

Tables vs. relation: Each table is a relation. Each row is a tuple of the relation. For example, the first row of the above table can be viewed as the atom STUDENT(111111111, John Doe, 123 Main St., Freshman).

Operations on tables. Needed since we would like to search for and/or combine information from tables, which might be huge. E.g., searching for something in a table, joining different tables to create new one, updating certain row, etc.

Several mathematical operations on relations can be used to define operations on tables.

Commercial DBMS helps make things simpler: it provides a convenient way for us to execute different operation on the tables; we only need to express what we want (as a *query*) and the DBMS will compute the result for us, i.e., we do not worry about *how can it be done*. Better DBMS is often equipped with more efficient engine to answer our queries.

SQL. A language that DBMS supports for users to query databases.

SQL: Basic SELECT statement. Syntax: SELECT FROM WHERE

Example:

```
SELECT Name FROM STUDENT WHERE Id = '111111111'
SELECT Id, Name FROM STUDENT WHERE Status = 'Senior'
SELECT * FROM STUDENT WHERE Status = 'Senior'
SELECT COUNT(*) FROM STUDENT WHERE Status = 'Senior'
```

Multi-table SELECT statement. Extended syntax: FROM's part includes a list of tables and WHERE's part is an expression over the attributes in the tables.

Suppose that grades are stored in a table named TRANSCRIPT with the attributes StudId, CrsCode, Semester, Grade.

```
SELECT Name, CrsCode, Grade
FROM STUDENT, TRANSCRIPT
WHERE StudId = Id and Status = 'Senior'
```

Note: See the difference in the condition 'WHERE'; the first equality combine the two tables in a certain way; the second one plays the role of a filter.

Query optimization. There are different ways for expressing a query. DBMS implements one way and often tries to optimize (the time) for computing the answers. However, we should learn to know how things can be done fast. Think about sorting!

Changing the content of a table. SQL provides commands for updating, inserting, and deleting a rows, a set of rows, or a table.

```
UPDATE STUDENT SET Status = 'Sophomore' WHERE Id = '111111111'
INSERT INTO STUDENT(Id, Name, Address, Status)
VALUES('999000999', 'Athur Ni', '123 Halloween', 'Freshman')
DELETE FROM STUDENT WHERE Id = '111111111'
```

Creating tables and specifying constraints. SQL allows us to create tables.

```
CREATE TABLE      STUDENT(
  Id                INTEGER,
  Name              CHAR(20),
  Address           CHAR(50),
  Status            CHAR(10),
  PRIMARY KEY(Id))
```

3 ACID Properties of Transaction

Transactions are programs that interact with databases (through the DBMS) to maintain the *correctness* of the databases. It is easier to think about the ACID properties if you think about your bank account and the database (of bank accounts) that your bank has.

Consistency. Executing a transaction at a state in which the database is consistent will result in a consistent database.

Withdrawing \$100 from an ATM does not only reduce your account's balance but also reduces the total amount of cash available at the ATM. And, you cannot get the \$100 if the balance is less than \$100.

Atomicity. The system must ensure that the execution of a transaction is completed as a whole or none.

If you withdraw \$100 and get the money, the balance must be reduced, the total amount of cash must be reduced. If you withdraw \$100 and did not get the money, nothing should change.

Durability. If a transaction is committed, its effects must remain throughout the lifetime of the database.

If you withdraw \$100, get the money, and your balance is 0, then it will stay so until you deposit some money.

Isolation. The overall effect of the execution of several transactions must be the same as if they are executed sequentially.

Withdrawing \$100 on two ATM simultaneously does not make you richer!

Why ACID? Database has consistency constraints or integrity constraints (IC); ICs are the characteristics of the information and must be maintained **at all time**.

For example, each student has a unique Id; a student cannot take two different courses at the same time; grades are letters 'A', 'B', 'C', 'D', 'F', 'I', 'S', or 'U'; the GPA of a student is the average of his grades; etc.

Transactions must ensure that they do not leave the database in a state violating the ICs of the system. Violating one of the ACID properties will lead to the violation of the ICs. Thus the requirement!

4 Exercises

- 2.1

MARRIED(Husband, Wife)	=	$\{(a, b) \mid a \text{ is husband of } b\}$
BROTHER(Brother, Person)	=	$\{(c, d) \mid c \text{ is brother of } d\}$
SIBLING(Person_1, Person_2)	=	$\{(e, f) \mid e \text{ is a sibling of } f\}$
BROTHER_IN_LAW	=	$\{(c, b) \mid \exists a \text{ such that MARRIED}(a, b) \text{ and BROTHER}(c, a)\} \cup$ $\{(c, a) \mid \exists b \text{ such that MARRIED}(a, b) \text{ and BROTHER}(c, b)\} \cup$ $\{(a, c) \mid \exists b \text{ such that MARRIED}(a, b) \text{ and SIBLING}(c, b)\}$

- 2.2

COURSESREGISTERFOR(Id, CrsCode)

COURSESTAKEN(Id, CrsCode, Grade)

- 2.4

```
CREATE TABLE TRANSCRIPT(  
  StudId      INTEGER,  
  CrsCode     CHAR(6),  
  Semester    CHAR(6),  
  Grade       CHAR(1),  
  PRIMARY KEY (StudId, CrsCode, Semester))
```