

# CS471, Programming Language Structure I

## Sample answer to Assignment 3, Fall 2002

The syntax is altered with the addition of a multiple assignment statement. This can be done recursively:

```
S ::= Ilist = Elist | ...
Ilist ::= I | I, Ilist
Elist ::= E | E, Elist
```

There is no way in BNF to specify an equal number of Is and Es.

There are no new domains needed, except that there must be operations for processing a list. These are head and tail:

```
head(I, Ilist) = I
tail(I, Ilist) = Ilist
```

the new valuation function for the multiple assignment is then:

$$M[[Ilist = Elist]] s = \text{if tail}(Ilist) \text{ is empty then } M[[\text{head}(Ilist) = \text{head}(Elist)]] s \\ \text{else } M[[\text{tail}(Ilist) = \text{tail}(Elist)]] (M[[\text{head}(Ilist) = \text{head}(Elist)]] s)$$

Note that this executes the tails of the lists in the store returned by executing the head. Thus the assignment:

$$x, y = 4, x$$

will first change  $x$  to 4, and then change  $y$  to 4, the new value of  $x$ . It also assumes that there are an equal number of Is and Es in the two lists. To do “parallel” assignment, in which the store does not change for each assignment:

$$M[[Ilist = Elist]] s = M[[Ilist = Elist]] s \\ M[[Ilist = Elist]] s_0 s_1 = \text{if } Ilist \text{ is empty then } s_1 \\ \text{else } M[[\text{tail}(Ilist) = \text{tail}(Elist)]] s_0 (M[[\text{head}(Ilist) = \text{head}(Elist)]] s_0)$$

In this version, the original store,  $s_0$  remains fixed, but the changes are accumulated in a second store,  $s_1$ .

Either version is a correct answer to the problem.