

? - cousin (mary, jake).

$\left\{ \begin{array}{l} X_1 = \text{mary} \\ Y_1 = \text{jake} \end{array} \right.$

parent( $P1_1$ , mary), parent( $P2_1$ , jake), sibling( $P1_1$ ,  $P2_1$ )

$\left| P1_1 = \text{john} \right.$

parent( $P2_1$ , jake), sibling(john,  $P2_1$ )

$\left| P2_1 = \text{susan} \right.$

sibling(john, susan)

$\left| \begin{array}{l} \text{mother}(M_2, \text{john}), \text{mother}(M_2, \text{susan}) \end{array} \right.$

$\left| \begin{array}{l} \text{parent}(M_2, \text{john}), \text{female}(M_2), \text{mother}(M_2, \text{susan}) \end{array} \right.$

$\left| M_2 = \text{jim} \right.$

female(jim), mother(jim, susan)



|  
father( $F_3$ , john), father( $F_3$ , susan)  
|  
parent( $F_3$ , john), male( $F_3$ ), father( $F_3$ , susan)  
|  $F_3 = \text{john}$   
male(john), father(john, susan)  
|  
father(john, susan)  
|  
parent(john, susan), male(john)  
|  
male(john)  
|  
yes

? - ancestor(jim, mary)

|  $X_1 = jim$   
|  $Y_1 = mary$

parent(jim, mary)

|  
fail, backtrack

|  $X_2 = jim$   
|  $Y_2 = mary$

parent(jim,  $Z_2$ ), ancestor( $Z_2$ ,  
mary)

|  $Z_2 = john$

ancestor(john, mary)

|  $X_3 = john$

|  $Y_3 = mary$

parent(john, mary)

|  
yes



a :- b, c, d.

a :- e, f.

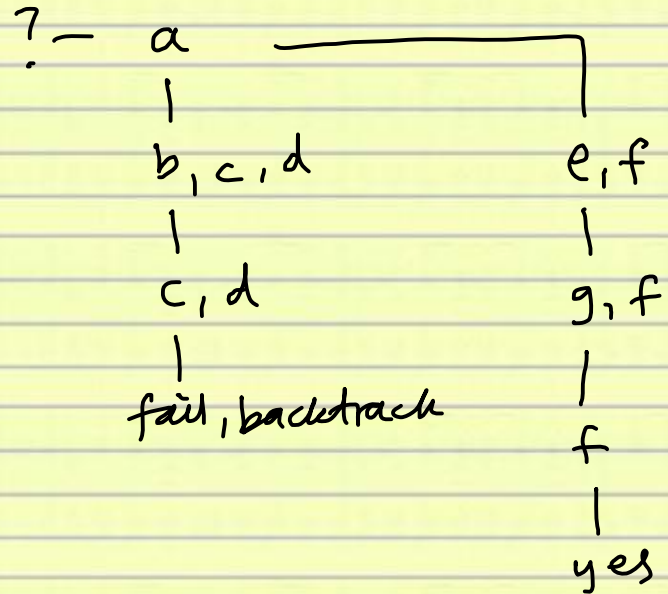
e :- g.

b.

g.

f.

d.



Call/Exit/Redo/Fail model

?- a

CALL a

CALL b

EXIT b

CALL c

FAIL c

REDO a

CALL e

CALL g

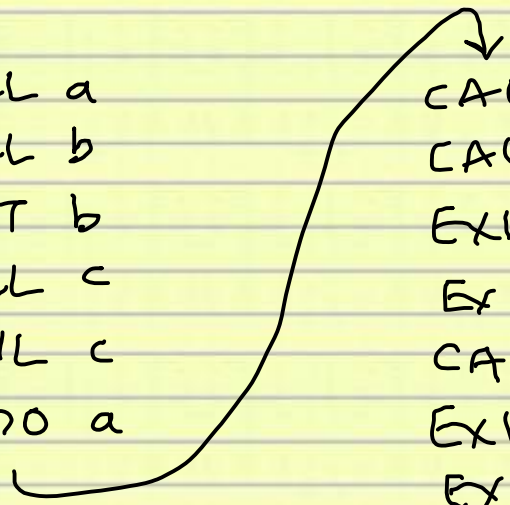
EXIT g

EXIT e

CALL f

EXIT f

EXIT a



$a :- b, c, d.$

$a :- e, f.$

$e :- g.$

$g.$

$b.$

$f.$

$d.$

Goal 'stack' - initially just a

a

↑

a b c d

↑

a b c d

↑

a

↑

a e f

↑

a e g f

↑

a e g f

↑

a e g

d

yes

Unification needs to very general

- matching any two terms, with variables (possibly)

terms:  $a(x)$

$a$

$a(b(x))$

$a(b(x), c(y))$

$a([x|y]) \Rightarrow a(\cdot(x, y))$

These are stored as nested structures

$a(b(c(x)))$  is represented as

Term a

Term b

Term c

Var X

e.g.  $a(b(x))$  matches  $a(b(c))$  with X bound to c

We need to record bindings so they can be undone on backtracking



We have looked at two semantics for Prolog

1. procedural (match, backtrack etc.)
2. declarative (logical statements)

There are 2 extra-logical features of Prolog

1. cut
2. negation

Cut - procedural mechanism for cutting off backtracking





Example using cut: member

member( $x, [x|_]$ ).

member( $x, [_|T]$ ) :- member( $x, T$ ).

- assume that a goal using member succeeds
- assume a later goal fails. Having succeeded, we matched the fact. However anything that matches the fact also matches the recursive rule. Backtracking to member causes (useless) search through the rest of the list.

? - member( $a, [b, a, c]$ ), fail.

$x_1 = a$   
 $T_1 = [a, c]$

member( $a, [a, c]$ ), fail

$x_2 = a$

fail, backtrack

$x_3 = a$

$T_3 = [c]$

member( $a, [c]$ )

|

