

append

e.g. (append '(a b) '(c d)) => (a b c d)

(define (append l1 l2)

(if (null? l1)

l2

(cons (car l1) (append (cdr l1) l2))))

e.g. (append '(a b) '(c d))

= (cons 'a (append '(b) '(c d)))

= (cons 'a (cons 'b (append () '(c d)))))

= (cons 'a (cons 'b '(c d)))

= (a b c d)

tail-recursive version

(define (append l1 l2))
 (define (append1 l1 l2))

(if (null? l1)
 L2

(append1 (cdr l1) (cons (car l1) l2))))

(append1 (reverse l1) l2))

assume (reverse '(a b c)) => (< b a)

e.g. (append '(a b) '(c d))

= (append1 '(b a) '(c d))

= (append1 '(a) '(b c d))

= (append1 () '(a b c d))

= (a b c d)

tail-recursive reverse

(define (reverse L))

(define (reverse1 L1 L2))

(if (null? L1)
 L2

(reverse1 (cdr L1) (cons (car L1) L2))))

(reverse L ()))

e.g. (reverse '(a b c))

= (reverse1 '(a b c) ())

= (reverse1 '(b c) '(a))

= (reverse1 '(c) '(b a))

= (reverse1 () '(c b a))

= (c b a)

Functions as first-class values

e.g. $(\text{map addOne } '(1 2 3)) \Rightarrow (2 3 4)$

$(\text{define (addOne } n) (+ n 1))$

$(\text{define (map } f L))$

$(\text{define (map1 } f L1 L2))$

$((\text{if (null? } L1)$

$L2$

$(\text{map1 } f (\text{cdr } L1))$

$(\text{cons } (f (\text{car } L1)) L2))))$

$(\text{map1 } f (\text{reverse } L1) ()))$

e.g. $(\text{map addOne } '(1 2 3))$

$= (\text{map1 addOne } '(3 2 1) ())$

$= (\text{map1 addOne } '(2 1) (\text{cons } (\text{addOne } 3) ()))$

$= (\text{map1 addOne } '(1) (\text{cons } (\text{addOne } 2) '()))$

$= (\text{map1 addOne } () (\text{cons } (\text{addOne } 1) '()))$

$\cdot (2 3 4)$

Note that map expects a function of one parameter as its argument. If it doesn't then an error will occur.

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(define (const1 a) (cons a ()))  
e.g. (map const1 '(1 a 2))  
      => ((1) (a) (2))
```

functions that take other functions as parameters are sometimes called "functionals"

filter - filters out items from a list according to some function

e.g. (filter odd? '(1 2 3 4)) => (2 4)

(define (odd? n) (= (mod n 2) 1))

(define (filter f L))

(cond ((null? L) ()))

((f (car L)) (filter f (cdr L))))

(else (cons (car L) (filter f (cdr L))))))

e.g. (filter odd? '(1 2 3 4))

= (filter odd? '(2 3 4))

= (cons 2 (filter odd? '(3 4)))

= (cons 2 (filter odd? '(4)))

= (cons 2 (cons 4 (filter odd? '()))))

= (cons 2 (cons 4 ())) = (2 4)

(filter const '(1 2 a))

Since const always returns a value, it is
the same as returning #t (like C, where a
non-zero value counts as true)

There is a "zero" value, called nil - this can
serve as false (#f) if returned by a function
and used in a test expression.

reduce - reduces a list to a single value using a binary function

e.g. $(\text{reduce} + '(1 2 3 4) \emptyset) \Rightarrow 10$

$(\text{define} (\text{reduce} f L n)$

$(\text{if} (\text{null? } L)$

n

$(f (\text{car } L) (\text{reduce} f (\text{cdr } L) n)))$

e.g. $(\text{reduce} + '(1 2 3) \emptyset)$

$= (+ 1 (\text{reduce} + '(2 3) \emptyset))$

$= (+ 1 (+ 2 (\text{reduce} + '(3) \emptyset)))$

$= (+ 1 (+ 2 (+ 3 (\text{reduce} + () \emptyset))))$

$= (+ 1 (+ 2 (+ 3 \emptyset)))$

$= 6$

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(reduce * '(1 2 3) \emptyset) => \emptyset

(reduce * '(1 2 3) 1) => 6

(reduce cons '(a b c) ()) => (a b c)

(This copies the list)

Semantics of Scheme

- could write an interpreter
- could use denotational semantics
- instead, we will look at a meta-circular interpreter — an interpreter for Scheme, written in Scheme

The meta-circular interpreter is an expression evaluator. In Scheme, this is called eval, and uses the fact that expressions and lists have the same syntax.

e.g. (define x '(+ 1 2))

x => (+ 1 2)

(eval x) => 3

(define x '(a b c))

(eval x) => error

Overview of meta-circular interpreter :

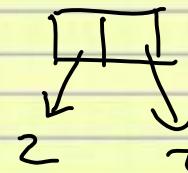
- use a binding environment for parameters
- represent it as a list

e.g. $((x.1) (x.(a\ b\ c)) (z.a))$

each item in the list is a binding pair,
the LHS is a parameter name, the RHS is the
value the parameter is bound to.

This is basically a cons cell

e.g. $(\text{cons}\ z\ 3) \Rightarrow (z.3)$



The list above is

