Detecting Patterns in Traces of Procedure Calls

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Objective

- Execution traces are important to understand the *behavior* and sometimes the *structure* of a software system
- Execution patterns can bridge the gap between *low level system* components and high level domain concepts
 - And hence help program comprehension
- We need:
 - an efficient way for *detecting* these patterns.
 - to understand when two patterns can be considered as equivalent
- We focus on patterns of procedure calls

Why Traces of Procedure Calls?

- The abstraction level of traces of procedure calls seems well suited for program comprehension.
 Applies to methods in OO as well
- Alternatives:
 - Inter-process messaging
 - Not a tree structure
 - Statement level
 - Vastly more detail

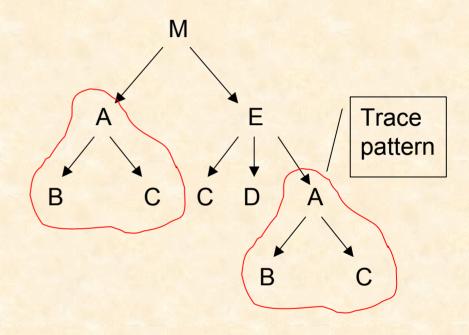
Definition of a Trace Pattern

"A sequence of calls that occurs <u>repetitively</u> but <u>non-contiguously</u> in several places in the trace"

- Zayour and Lethbridge
- We add: instances of a given pattern do not need to be identical
- Ideally, a trace pattern corresponds to an abstract domain concept.
 - E.g. a user identifiable aspect of some feature
 - But reality is far from the ideal

Trace Structure

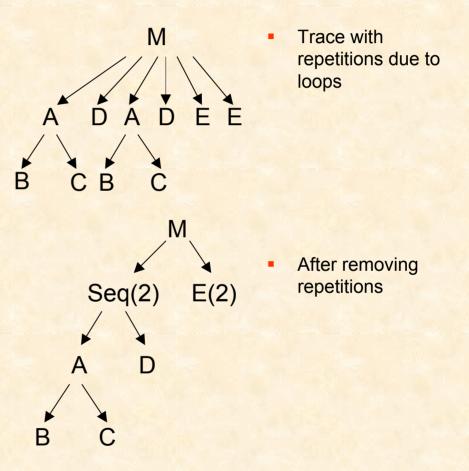
- A trace of procedure calls from a single thread is represented by a rooted labeled ordered tree
- A trace pattern is represented as a non-contiguous repeated subtree.



Tree representation of the trace:
 M(A(B, C), E(C, D, A(B, C)))

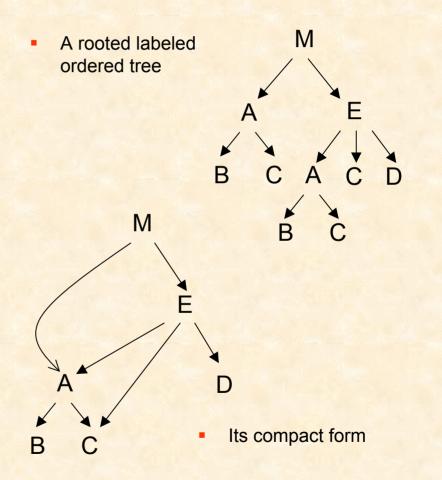
Trace Preprocessing

- Sequences of calls due to loops and recursion encumber the trace
- They need to be removed
- The tree structure is maintained by adding virtual calls



Pattern Extraction

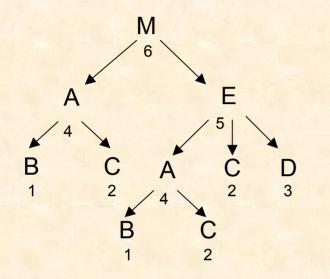
- Based on Valiente's algorithm for the common subexpression problem
- (The problem was introduced by J.P. Downey, R. Sethi and R.E. Tarjan)
- Any rooted tree can be transformed into its most compact form by representing common subtrees only once



Valiente's Algorithm

- Traverse the tree in a bottom-up fashion
- Assign a positive integer between 1 and n In represents the size of the tree
- Two nodes *n1* and *n2* are assigned the same *certificate* if the trees rooted at them are similar according to *predefined matching criteria*
- To compute the certificate, each node is assigned a signature
- The signature of a node *n* consists of its label and the certificates of its direct children, if there are any.
 - This makes it unique
- A global hash table is used to store the certificates and signatures and ensure that similar subtrees will always hash to the same element.
- Our contribution: Examine the matching criteria

Example



 If we consider exact match only, identical subtrees have the same certificates

Certificate	Frequency	Signature
1	2	В
2	3	С
3	1	D
4	2	A 1 2
5	1	E 4 2 3
6	1	M 4 5

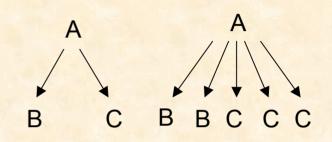
- Global table that corresponds to this tree.
- We mainly traverse the table and extract the subtrees (excluding the leaves) that have more than 1 incoming edge.
- E.g. 'A 1 2' is a pattern

Equivalent Patterns

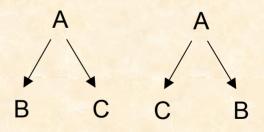
- Pattern matching criteria determine when two sequences of calls can be considered equivalent
 - Will effect usefulness and performance
- Users can select the criteria according to their knowledge of the system
 - E.g. identical patterns might be useful to novices but less useful to experts.
- De Pauw, Lorenz, Vlissides and Wegman suggested a list of matching criteria that are used for OO systems.
 Some of them are also useful for procedural systems.

De Pauw et al. Pattern Matching Criteria Useful For Procedural Systems

- **Identity**: Two sequences of calls are similar if they have the same topology:
 - Same call structure
 - Same order of calls
 - D Etc.
- Ignoring Repetition: repetition due to loops and recursions can be ignored when looking for patterns
- **Ignoring Ordering**: Order of calls might not be important at some levels of the call tree.
- **Depth limited**: allows comparing two subtrees up to a certain depth.
 - Deeper calls ignored



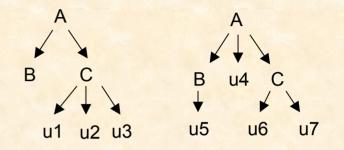
These two sequences are similar if repetitions are ignored



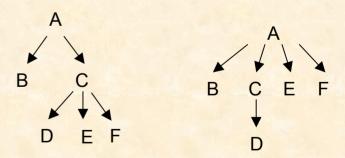
 Two sequences are similar if the order is ignored

New Matching Criteria

- Removing Utilities: Requires
 automatic detection of utility routines
- **Distance Limited**: Trees within a certain edit distance are considerd the same.
- Flattening: Ignore structure of tree
 - consider only routines encountered
 - might be useful for experts who are not interested in the call structure



 These patterns are considered equivalent if utility routines (u_i) are ignored



 The edit distance can reveal that these two sequences can be considered as the same pattern

Conclusions and Future Work

- We need to validate the matching criteria and analyze at which level of the tree they can be applied usefully
- We also need to study how they can be combined.
- It is also important to understand the relationship between the use of the matching criteria and the user's knowledge of the system
- Goal: automatic detection of patterns that most likely correspond to high level concepts.

Why this work will fail!

- We might have trouble finding ways for the criteria to work together
- There might be no intuitive meaning attributable to edit distance in the general context
- We might not find a universally appropriate definition of a 'utility'

Why this work will succeed!

- We have already done some experiments in compressing traces
- High interest from SEs in browsing smaller traces
- We have been told by SEs that our proposed matching criteria correspond to what people want