

December 2018 Discrete Mathematics Qualifying Exam

Closed book closed notes

1. (25 points) Given the following compound proposition (expression in propositional logic), find a logically equivalent expression that uses only \wedge and \neg and that has no more than one occurrence of p , no more than one occurrence of q , and no more than one occurrence of r):

$$(r \rightarrow (q \leftrightarrow r)) \rightarrow p$$

Solution:

It could be done using truth tables or using logical equivalences. The following solution uses logical equivalences. (Note: you do not have to name the laws you are using as long as they are used correctly)

$$\begin{aligned} & (r \rightarrow (q \leftrightarrow r)) \rightarrow p \\ & \quad \text{(eliminate biconditional using Conditional identities)} \\ \equiv & (r \rightarrow ((q \rightarrow r) \wedge (r \rightarrow q))) \rightarrow p \\ & \quad \text{(eliminate implications using Conditional identities)} \\ \equiv & (\neg(\neg r \vee ((\neg q \vee r) \wedge (\neg r \vee q)))) \vee p \\ & \quad \text{(move negation inward using De Morgan's laws)} \\ \equiv & (\neg \neg r \wedge \neg((\neg q \vee r) \wedge (\neg r \vee q))) \vee p \\ & \quad \text{(use Double negation law)} \\ \equiv & (r \wedge \neg((\neg q \vee r) \wedge (\neg r \vee q))) \vee p \\ & \quad \text{(move negation inward using De Morgan's laws)} \\ \equiv & (r \wedge (\neg(\neg q \vee r) \vee \neg(\neg r \vee q))) \vee p \\ & \quad \text{(move negation inward using De Morgan's laws)} \\ \equiv & (r \wedge ((\neg \neg q \wedge \neg r) \vee (\neg \neg r \wedge \neg q))) \vee p \\ & \quad \text{(use Double negation law)} \\ \equiv & (r \wedge ((q \wedge \neg r) \vee (r \wedge \neg q))) \vee p \\ & \quad \text{(use Distributive law)} \\ \equiv & ((r \wedge (q \wedge \neg r)) \vee (r \wedge (r \wedge \neg q))) \vee p \\ & \quad \text{(use Commutative and Associative laws)} \\ \equiv & (((r \wedge \neg r) \wedge q) \vee ((r \wedge r) \wedge \neg q)) \vee p \\ & \quad \text{(use Complement and Idempotent laws)} \\ \equiv & ((F \wedge q) \vee (r \wedge \neg q)) \vee p \\ & \quad \text{(use Commutative and Domination laws)} \\ \equiv & (F \vee (r \wedge \neg q)) \vee p \\ & \quad \text{(use Commutative and Identity laws)} \\ \equiv & (r \wedge \neg q) \vee p \\ & \quad \text{(use De Morgan and Double negation laws to remove } \vee) \\ \equiv & \neg(\neg(r \wedge \neg q) \wedge \neg p) \end{aligned}$$

2. (25 points) Determine whether the following statement is true or false. Justify your answer with a proof or a counterexample, as appropriate.

For all real numbers x and y , if $-y^3 + x^2y + 2y + 1 \leq x^3 + xy^2 + 2x$, then $y \leq x$.

Solution: The statement is false. Counterexample: Take $x=0$ and $y=2$. Then,

$$\text{LHS} = -y^3 + x^2y + 2y + 1 = -8 + 0 + 4 + 1 = -3.$$

$$\text{RHS} = x^3 + xy^2 + 2x = 0.$$

$$-3 \leq 0 \text{ However, } y \not\leq x.$$

3. (25 points) Consider $\{b_n\}$ defined by $b_1 = 1$, $b_2 = 5$, and $b_n = b_{n-1} + 2b_{n-2}$ for all integers $n \geq 3$. Prove that for all integers $n \geq 1$, $b_n = 2^n + (-1)^n$.

Proof by strong induction.

Base cases.

$$\text{When } n = 1, \text{ we have } b_1 = 1 \text{ and } 2^1 + (-1)^1 = 2 - 1 = 1.$$

$$\text{When } n = 2, \text{ we have } b_2 = 5 \text{ and } 2^2 + (-1)^2 = 4 + 1 = 5.$$

Inductive step.

Assume that $b_n = 2^n + (-1)^n$ for all integers n from 1 to k for some integer $k \geq 2$ (strong inductive hypothesis). We need to show that $b_{k+1} = 2^{k+1} + (-1)^{k+1}$.

Since $k \geq 2$, we have that $k+1 \geq 3$ and $(k+1) - 2 \geq 1$.

By definition,

$$\begin{aligned} b_{k+1} &= b_k + 2b_{k-1} \\ &= 2^k + (-1)^k + 2(2^{k-1} + (-1)^{k-1}) && \text{by strong inductive hypothesis} \\ &= 2^k + (-1)^k + 2^k + 2(-1)^{k-1} \\ &= 2 \cdot 2^k + (-1)(-1)^{k-1} + 2(-1)^{k-1} \\ &= 2^{k+1} + (-1)^{k-1} \\ &= 2^{k+1} + (-1)^{k+1} \end{aligned}$$

This is what we needed to show.

3. Consider a set of strings of length n with alphabet $\{a, b, c, d\}$. This is a set of strings of length n , where each of the n symbols is one of $\{a, b, c, d\}$. n is a positive integer.
- (3 points) How many strings of length n with alphabet $\{a, b, c, d\}$ are there?
 - (3 points) How many of strings of length n do not contain a?
 - (3 points) How many of strings of length n contain at least one a?
 - (3 points) How many of strings of length n do not contain a and do not contain b?
 - (13 points) Calculate the number of strings of length n with alphabet $\{a, b, c, d\}$ such that each of a, b, and c occurs at least once.
For instance, if $n=3$, then the number of strings of length 3 with alphabet $\{a, b, c, d\}$ such that each of a, b, and c occurs at least once is 6 (which is the number of permutations of a, b, and c).
If $n=1$ or $n=2$, then the number of such strings is 0.
You need to derive a general formula for the number of strings of length n with alphabet $\{a, b, c, d\}$ that contain each of a, b, and c at least once.

Solution:

- a) Total number of n-strings of 4 symbols is 4^n .
- b) The number of n-strings that do not contain a is 3^n .
- c) The number of n-strings that contain at least one a is $4^n - 3^n$.
- d) The number of n-strings that do not contain a and do not contain b is 2^n .
- e) Let A be a set of n-strings that contain a. Let B be a set of n-strings that contain b. Let C be a set of n-strings that contain c. Then, \bar{A} is a set of n-strings that do not contain a. \bar{B} is a set of n-strings that do not contain b. \bar{C} is a set of n-strings that do not contain c.

The question is asking for $|A \cap B \cap C|$.

$$A \cap B \cap C = \overline{\bar{A} \cup \bar{B} \cup \bar{C}}. \text{ Therefore, } |A \cap B \cap C| = 4^n - |\bar{A} \cup \bar{B} \cup \bar{C}|.$$

By the inclusion-exclusion principle, $|\bar{A} \cup \bar{B} \cup \bar{C}| = |\bar{A}| + |\bar{B}| + |\bar{C}| - |\bar{A} \cap \bar{B}| - |\bar{A} \cap \bar{C}| - |\bar{B} \cap \bar{C}| + |\bar{A} \cap \bar{B} \cap \bar{C}|$.

$$|\bar{A}| = |\bar{B}| = |\bar{C}| = 3^n.$$

$$|\bar{A} \cap \bar{B}| = |\bar{A} \cap \bar{C}| = |\bar{B} \cap \bar{C}| = 2^n.$$

$$|\bar{A} \cap \bar{B} \cap \bar{C}| = 1.$$

Therefore,

$$|\bar{A} \cup \bar{B} \cup \bar{C}| = 3^n + 3^n + 3^n - 2^n - 2^n - 2^n + 1 = 3^{n+1} - 3 \cdot 2^n + 1.$$

The number of n-strings with alphabet $\{a,b,c,d\}$ such that each of a, b, and c occurs at least once is

$$|A \cap B \cap C| = 4^n - |\bar{A} \cup \bar{B} \cup \bar{C}| = 4^n - 3^{n+1} + 3 \cdot 2^n - 1.$$