
A-POL

a Partial Order Programming extension for Answer Set Programming

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PRESENTATION PLAN

- Preliminaries
- A-POL Description
- Implementation
- Examples and results

PREELIMINARIES

- Partial Order Programming
- Form of partial-order clauses:
 - $f(\text{terms}) \geq \text{expression} : - p_1(\text{terms}_1), \dots, p_n(\text{terms}_n).$
 - $f(\text{terms}) \leq \text{expression} : - p_1(\text{terms}_1), \dots, p_n(\text{terms}_n).$
- Partial Order Clauses as normal programs

A-POL DESCRIPTION

- What is A-POL?
- Main advantages:
 - Efficient
 - Domain size
- Main features:
 - Witness recovery
 - User defined partial orders
 - User defined libraries

A-POL DESCRIPTION

Shortest distance between two nodes of a graph

```
node(a).  node(b).  node(c).  node(d).
```

```
edge(a,b,1).  edge(b,c,2).  edge(a,d,4).
```

```
start(a).  end(c).
```

```
short(X)<= 0 :- start(X).
```

```
short(X)<= add(short(Y),C) :- edge(Y,X,C).
```

A-POL DESCRIPTION

A-POL accepts disjunctive clauses of the form:

$$p_1(t_1) \vee \dots \vee p_n(t_n) : -$$

$$b_1(t_{n+1}), \dots, b_m(t_{n+m}),$$

$$f_1(t_{n+m+1}) = r_1, \dots, f_s(t_{n+m+s}) = r_s.$$

and partial-order clauses of the form:

$$f(\text{terms}) \geq \text{expression} : -$$

$$p_1(\text{terms}_1), \dots, p_n(\text{terms}_n). / \text{info}(\text{iterms}).$$

$$f(\text{terms}) \leq \text{expression} : -$$

$$p_1(\text{terms}_1), \dots, p_n(\text{terms}_n). / \text{info}(\text{iterms}).$$

IMPLEMENTATION

Main Computation Cycle:

```
A-POL( $P$ ) {  
    order = getLevel( $P$ )  
    models =  $\emptyset$   
    for each  $cl \in$  order  
        if ( $cl$  are disjunctive clauses)  
            models = ASP(models, $cl$ )  
        else  
            models = POP(models, $cl$ )  
    return models  
}
```

IMPLEMENTATION

- Evaluation path for a stratified program:

$$\mathit{level}(cl_1) = 1$$

$$\mathit{level}(cl_2) = 2$$

$$\mathit{level}(cl_3) = 3$$

....

$$\mathit{level}(cl_n) = n$$

- ASP - we use DLV^a
- POP - general dynamic programming algorithm.

^a<http://www.dbai.tuwien.ac.at/proj/dlv/>

EXAMPLES AND RESULTS

Shortest Path

```
declare < short(node).  
  
node(a).  node(b).  node(c).  node(d).  
  
edge(a,b,1).  edge(b,c,2).  edge(a,d,4).  
  
start(a).  end(c).  
  
short(X)<= 0 :- start(X).  
  
short(X)<= add(short(Y),C) :- edge(Y,X,C). / way(Y,X).  
  
select way where short(X) <- end(X).
```

EXAMPLES AND RESULTS

DLV Shortest Path

```

way(X,Y,C) v other(X,Y,C) :- edge(X,Y,C).

:- edge(X,A,C), start(A), way(X, A,C).

:- edge(D,X,C), end(D), way(D,X,C).

:- edge(X,Y,C), edge(X,Y1,C1), way(X,Y,C),
   way(X,Y1,C1), Y != Y1.

:- edge(X,Y,C), edge(X1,Y,C1), way(X,Y,C),
   way(X1,Y,C1), X != X1.

r(X) :- start(X).

r(X) :- r(Y), edge(Y,X,C), way(Y,X,C).

:- end(D), not r(D).

: way(X,Y,C). [C:1]

```

EXAMPLES AND RESULTS

Experimental results for shortest path implementations

| # nodes | DLV | A-POL |
|----------------|------------|--------------|
| 10 | 0.16s | 0.65s |
| 20 | 4.77s | 0.84s |
| 40 | 4m 21.45s | 1.15s |
| 60 | 22m 57.02s | 2.21s |
| 80 | - | 2.85s |

EXAMPLES AND RESULTS

0-1 Knapsack

```

maxint = 20.

declare > ks(object,tweight).

declare > profit(object).

select take1 where ks(4,10).

object(0..4).  tweight(0..10).

weight(1,2).  weight(2,3).  weight(3,5).  weight(4,8).

profit(0)>= 0.  profit(1)>= 1.  profit(2)>= 2.

profit(3)>= 3.  profit(4)>= 5.

ks(N,K)>= ks(sub(N,1),K) :- le(1,N). / take1(N,K).

ks(N,K)>= add(profit(N),ks(sub(N,1),sub(K,X))) :-
    weight(N,X), le(X,K), le(1,N). / take1(N,K).

take(X1,Y) :- take1(X1,Y), not take(X2,Y), X2 > X1,
    object(X2), tweight(Y).

```

EXAMPLES AND RESULTS

Other problems coded in A-POL

- Optimal parentization for a matrix-chain multiplication
- Data-flow analysis equations
- ACM International Collegiate Programming Contest problems (fastfood, matrix multiplication)
- A simple capital budgeting problem^a
- Cartoon Co. example from the DLV manual (introducing aggregate functions)

^a<http://mat.gsia.cmu.edu/classes/dynamic/node2.html>

CONCLUSIONS AND FUTURE WORK

Conclusions

- Successfully implementation of A-POL
- Experimental results support our extension for ASP
- We present typical optimization problems coded in A-POL.

Future work

- Only a subset of DLV language is supported
- Explore the uses of partial-order clauses in other DLV front ends (like DLV-K).

THANK YOU!

<http://www.udlap.mx/is108990/apol>