

## Automata (Fall 2008) Qual Exam

Answer ALL questions

### Question 1

Given a DFA  $M = (Q, \Sigma, \delta, q_0, F)$ , we overload the function name  $\delta$  to define a function that maps letters in  $\Sigma$  to functions from  $Q$  to  $Q$ . Specifically, we define  $\delta : \Sigma \rightarrow (Q \rightarrow Q)$  such that  $\delta(a)(q) = q'$  iff  $\delta(q, a) = q'$  where  $a \in \Sigma$  and  $q, q' \in Q$ .

We can further extend  $\delta$  to define  $\delta : \Sigma^* \rightarrow (Q \rightarrow Q)$  inductively such that  $\delta(aw) = \delta(a) \circ \delta(w)$  where  $a \in \Sigma$ ,  $w \in \Sigma^*$  and  $\circ$  denotes function composition. That is,  $\delta(aw)(q) = \delta(a)(\delta(w)(q))$  for  $q \in Q$ . For the base case of the inductive definition,  $\delta(\epsilon)$  is defined to be the identity function such that  $\delta(\epsilon)(q) = q$  for  $q \in Q$ .

With respect to a DFA  $M = (Q, \Sigma, \delta, q_0, F)$ , we define an equivalence relation  $\equiv_M$  on strings in  $\Sigma^*$  such that  $w \equiv_M w'$  iff  $\delta(w) = \delta(w')$ .

With respect to a given language  $L$ , we define an equivalence relation  $\equiv_L$  on strings in  $\Sigma^*$  such that  $w \equiv_L w'$  iff  $\forall \alpha, \beta \in \Sigma^*$ ,  $\alpha w \beta \in L \iff \alpha w' \beta \in L$ .

In the following questions,  $M$  is assumed to be a DFA.

(a) Show that  $\equiv_M$  has finite index; that is, there are finitely many  $\equiv_M$  equivalence classes. You are required to give an upper bound on the number of equivalence classes in terms of  $n$ , which is the number of states in  $M$ .

(b) Show that  $w \equiv_M w'$  implies  $w \equiv_L w'$  where  $L$  is the language of  $M$ .

(c) Illustrate with an example to show that  $w \equiv_L w'$  does not necessarily imply  $w \equiv_M w'$  where  $L$  is the language of  $M$ .

Question 2

Let  $\Sigma$  be a finite alphabet. For any string  $w = a_1a_2 \dots a_n$ , where  $a_i \in \Sigma$ , the reverse of  $w$ , written  $w^R$ , is the string  $w$  in reverse order,  $a_n \dots a_2a_1$ . For any language  $L$ ,  $L^R = \{w^R \mid w \in L\}$ .

- (a) If  $L$  is regular, is  $L^R$  always regular? Justify your answer.
- (b) If  $L$  is context-free, is  $L^R$  always context-free? Justify your answer.
- (c) If  $L$  is deterministic context-free, is  $L^R$  always deterministic context-free? Justify your answer. (Note that a language is deterministic context-free if the language can be recognized by a deterministic pushdown automaton.)